## Lecture Hierarchical Planning

# Chapter: Heuristics for (Hierarchical) Planning Problems

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# ulm university universität **UUUIM**

ask Decomposition Graph

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Compilation Technique

Summary

### **Overview:**

- 1 Task Decomposition Graph
- 2 Landmarks

# 3 TDG-c & TDG-m

- Cost-sensitive TDG Heuristic, TDG-c
- Modification-sensitive TDG Heuristic, TDG-m
- Properties of TDG Heuristics
- TDG Recomputation

# 4 Compilation Technique



Introduction ●○	Task Decomposition Graph	TDG-c & TDG-m 00000000000000000	Compilation Technique	

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Introduction		TDG-c & TDG-m	
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- With the right combination of algorithm and heuristic, we can also provide optimality guarantees.
- In classical planning, heuristics estimate the number of actions (or their costs) that need to be applied to reach a goal state.
- In POCL planning, heuristics could also estimate the number of required modifications (which, in addition to task insertion, may estimate the number of ordering constraints and causal links that need to be added).



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What's the different to hierarchical planning?

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- We have to deal with abstract tasks (see above).
- Additional challenges for decomposition-based planning:
  - There is no current state, only the current partial plan.
  - Search nodes get bigger and bigger. Thus, paradoxically, the problem gets *harder* the closer we approach a solution.



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Introduct	tion			

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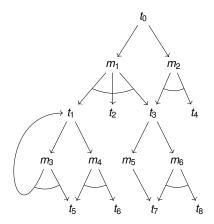
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- In contrast, the decomposition graph represents the planning domain (cf. introductory chapter).
- Due to possibly cyclic methods, DTs are in general no sub structures of TDGs.



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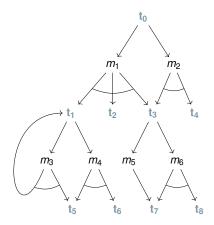
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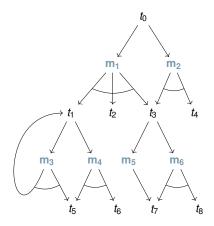


- A TDG is a bipartite graph  $\mathcal{G}$  $\langle N_T, N_M, E_{(T,M)}, E_{(M,T)} \rangle$  with
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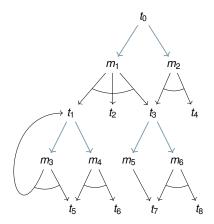


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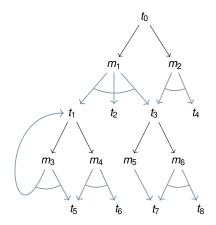


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Prerequisites of Formal Definition						

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  - VC<sub>m</sub> is also a set of variable constraints to relate the variables in  $\overline{\tau}$  to the variables in  $tn_m$ .
- Ground<sub>VC</sub>(tn) denotes the set of all possible groundings of tn by also taking into account the variable constraints VC.



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Let  $\mathcal{P} = (\mathcal{D}, s_I, c_I)$  with  $\mathcal{D} = (\mathcal{L}, \mathcal{P}, \delta, \mathcal{C}, \mathcal{M})$  be an HTN planning problem. The graph  $\mathcal{G} = \langle V_T, V_\mathcal{M}, E_{T \to \mathcal{M}}, E_{\mathcal{M} \to T} \rangle$  is called the TDG of  $\mathcal{P}$  if it holds:

**1 base case** (task vertex for the given task)  $c_l \in V_T$ , the TDG's root.



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**2** method vertices (derived from task vertices) Let  $v_t \in V_T$  with  $v_t = t(\bar{c})$  and  $(t(\bar{\tau}), tn_m, VC_m) \in M$ . Then, for all  $v_m \in Ground_{VC_m \cup \{\bar{\tau} = \bar{c}\}}(tn_m)$  holds: •  $v_m \in V_M$  •  $(v_t, v_m) \in E_{T \to M}$ .



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# 4 tightness

 ${\cal G}$  is minimal, such that 1. to 3. hold.



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- What's its construction time? That depends on the input:
  - Ground Model: Polynomial time.
  - Lifted Model: Exponential time (due to the grounding).



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  - Repeat until every compound task (i.e., their groundings) is in the TDG.



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- $\rightarrow$  More details on next slide.



Introduction	Task Decomposition Graph ○○○○○○●○○	TDG-c & TDG-m 000000000000000000000000000000000000	Compilation Technique	
Ontimiza	ations			

- If a problem instance is given, we can perform a reachability analysis and thereby reducing its size:
  - We can restrict the TDG to those tasks reachable from the initial task (network). (This follows already from the definition.)
  - $\rightarrow$  Top-down reachability analysis.
  - We can restrict the TDG's primitive tasks to those reachable from the initial (current?) state.
  - ightarrow Bottom-up reachability analysis.
  - We can restrict to compound tasks with at least one method.
  - We can restrict to compound tasks that allow a primitive decomposition.
  - We can restrict to task networks without "eliminated elements".
- $\rightarrow\,$  More details on next slide.

Note: *Technically*, any modification to the TDG will violate its definition. We still refer to the resulting structures as TDGs, though.



Introduction	Task Decomposition Graph	TDG-c & TDG-m 0000000000000000	Compilation Technique	

Step 1: Construct PG to find reachable ground primitive tasks.



Introduction 00	Task Decomposition Graph ○○○○○○○●○	TDG-c & TDG-m 00000000000000000	Compilation Technique	

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  - Repeat until nothing can be deleted.
- Step 5: Since the set of reachable primitive tasks may have changed, we can repeat all previous steps (possibly multiple times).This step does usually not pay off empirically.



Task Decomposition Graph	TDG-c & TDG-m	
00000000		

Restricting the TDG – Further Optimization

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  - We build a task parameter- and predicate parameter-free TDG.



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  - We build a task parameter- and predicate parameter-free TDG.
  - This TDG can be built very efficiently.



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  - We first perform a *parameter-relaxed* top-down analysis.
  - We build a task parameter- and predicate parameter-free TDG.
  - This TDG can be built very efficiently.
  - This overestimates the number of reachable tasks, but already rules out some unreachable actions for the PG construction.



	Task Decomposition Graph	Landmarks ●000000000	TDG-c & TDG-m 0000000000000000	Compilation Technique	
Introduction ar	nd Definitions				
Non-Hie	rarchical Landmarks				

The concept of *landmarks* originates from *classical* planning, but it was transferred to hierarchical planning later on.



Introduction	Task Decomposition Graph	Landmarks ●000000000	TDG-c & TDG-m 00000000000000000	Compilation Technique	
Introduction ar	nd Definitions				
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- What are *trivial* fact landmarks? The initial state and goal description.
- A classical action landmark is an action that is part of every solution.



Introduction	Task Decomposition Graph	Landmarks O●O0000000	TDG-c & TDG-m 00000000000000000	Compilation Technique	
Introduction ar	nd Definitions				
Hierarch	ical Landmarks				

- A hierarchical landmark is a task (primitive or compound) that occurs on any sequence of decompositions from the initial task (network) to any solution.
- A formal definition will be provided or has to be found in the exercises.



Introduction	Task Decomposition Graph	Landmarks 00●0000000	TDG-c & TDG-m 00000000000000000	Compilation Technique	
Introduction a	nd Definitions				
Exploitat	tion of Landmarks				

Why are we interested in landmarks?



Introduction	Task Decomposition Graph	Landmarks 00●0000000	TDG-c & TDG-m 00000000000000000	Compilation Technique	
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Introduction	Task Decomposition Graph	Landmarks 0000000000	TDG-c & TDG-m 0000000000000000	Compilation Technique	
Introduction ar	nd Definitions				
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Introduction	Task Decomposition Graph	Landmarks 00●0000000	TDG-c & TDG-m 0000000000000000	Compilation Technique	
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  - Explanations: Explanations basing on landmarks might be more convincing.
  - Heuristics: Later in this section!





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- → All in all: Reduction from the *Plan-Nonexistence decision problem* to the *Landmark decision problem*.



Introduction	Task Decomposition Graph	Landmarks ○○○○●○○○○○	TDG-c & TDG-m 0000000000000000	Compilation Technique	
Computational	Complexity				
Classical	Landmarks (Memb	ership)			

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  - Given a fact  $v \in V$ , remove all actions that generate v and check whether the problem is still solvable.
  - *v* is a landmark if and only if it has become unsolvable.





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- Hierarchical Landmarks (Hardness)
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    - The third method maps c\* to the empty task network thereby allowing a trivial solution.





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    - The third method maps c\* to the empty task network thereby allowing a trivial solution.
  - If  $c^* \notin C$  is a landmark, then the original task was unsolvable.
  - → All in all: Reduction from the *Plan-Nonexistence decision problem* to the *Landmark decision problem*.



Introduction	Task Decomposition Graph	Landmarks ○○○○○○●○○○	TDG-c & TDG-m 0000000000000000	Compilation Technique	
Computationa	I Complexity				

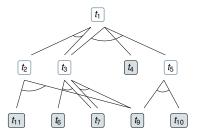
Hierarchical Landmarks (Membership)

- Is the problem semi-decidable?
- $\rightarrow$  Exercise!





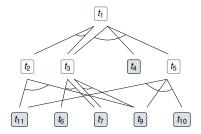
General idea: Compute the intersection of all partial plans that belong to the same compound task.





Introduction	Task Decomposition Graph	Landmarks ○○○○○○○○●○	TDG-c & TDG-m 00000000000000000	Compilation Technique	
Computation					
Landmar	k Computation, Imp	roved			

We can still do better than that, though...







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- → Both approaches lack from the problem that the landmarks as computed are those that will be used no matter what, i.e., one *cannot prevent* having them in a sequence of decompositions – limiting their usefulness.
  - We can use all primitive landmarks as the basis for state-based heuristics!
- → This allows to use *any* classical heuristic (or classical landmark technique!).



Introduction	Task Decomposition Graph	TDG-c & TDG-m ●o○○○○○○○○○○○○○○	Compilation Technique	
Introduction				
Introduct	ion			

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Introduction				
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- Even in these cases, the TDG holds valuable information that can be exploited to estimate goal distances...



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- But how..? We know that:
  - all tasks within a method need to be "accomplished" (applied or decomposed).
  - For each compound task, only one of its methods needs to be applied.



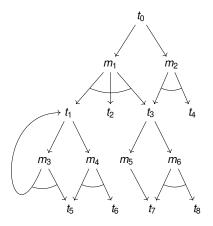
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- The landmark heuristics take only simple landmarks into account.
- More precisely: one can easily rewrite a model, such that none of its landmarks gets discovered.
- Even in these cases, the TDG holds valuable information that can be exploited to estimate goal distances...
- But how..? We know that:
  - all tasks within a method need to be "accomplished" (applied or decomposed).
  - For each compound task, only one of its methods needs to be applied.
- Note: For convenience, we later write  $t(\bar{\tau}) \in T$  as shorthand for  $t(\bar{\tau}) = \alpha(t'), t' \in T$ .



Introduction	Task Decomposition Graph	TDG-c & TDG-m ⊙●○○○○○○○○○○○○○○	Compilation Technique	
Introduction				

## Overview



Exploit TDG for effort estimation.

Step 1:

Compute the TDG.

## Step 2:

Compute TDG-based estimates  $h_T(t)/h_M(t)$  for each task/method node in the TDG (*once* via preprocessing).

## Step 3:

For search node (task network or partial plan) *tn* and its tasks *T*, compute h(tn) based on the estimates for the  $t \in T$ .

- Via estimating the costs of missing actions → TDG-c.
- Via estimating the still required modifications → TDG-m.



Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○●○○○○○○○○○○○○○	Compilation Technique	
Cost-sensitive	TDG Heuristic, TDG-c			
Motivatio	าท			

To obtain good (or cheap) solutions heuristically, we need to estimate the costs of a plan that can be developed from a current task network.



Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○●○○○○○○○○○○○○	Compilation Technique	
Cost-sensitive	e TDG Heuristic, TDG-c			
Motivotic	22			

- To obtain good (or cheap) solutions heuristically, we need to estimate the costs of a plan that can be developed from a current task network.
- We thus exploit the TDG and use its action costs as basis for estimates.



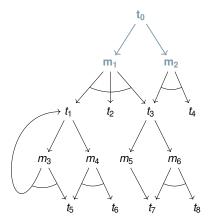
Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○●○○○○○○○○○○○○	Compilation Technique	
Cost-sensitive	TDG Heuristic, TDG-c			
Mativatio	22			

- To obtain good (or cheap) solutions heuristically, we need to estimate the costs of a plan that can be developed from a current task network.
- We thus exploit the TDG and use its action costs as basis for estimates.
- The resulting heuristic will be admissible (trivial).



Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○○●○○○○○○○○○○○○	Compilation Technique	
Cost-sensitive	TDG Heuristic, TDG-c			

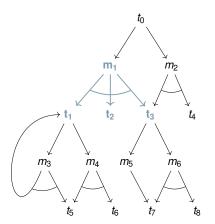
# Example:



$$h_T(t_0) = min\{h_M(m_1), h_M(m_2)\}$$

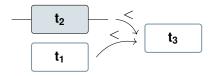


Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○○●○○○○○○○○○○○	Compilation Technique	
Cost-sensitive	TDG Heuristic, TDG-c			



# Example:

Method  $m_1 = (t_0, t_n)$  with task network *tn*:

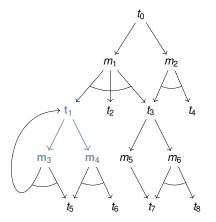


$$h_M(m_1) = \sum_{t_i \in \{t_1, t_2, t_3\}} h_T(t_i)$$



Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○○●○○○○○○○○○○○○	Compilation Technique	
Cost-sensitive	TDG Heuristic, TDG-c			

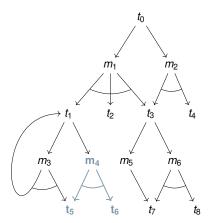
# Example:



$$h_T(t_1) = min \{h_M(m_3), h_M(m_4)\}$$

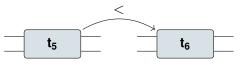


Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○○●○○○○○○○○○○○	Compilation Technique	
Cost-sensitive	TDG Heuristic, TDG-c			



Example:

Method  $m_4 = (t_1, t_n)$  with task network *tn*:



$$h(m_4) = c(t_5) + c(t_6)$$



Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○○○○○○○○○○○○○○	Compilation Technique	
Cost-sensitive	e TDG Heuristic, TDG-c			
TDG-c C	Computation			

# Let $\mathcal{G} = \langle V_T, V_M, E_{T \to M}, E_{M \to T} \rangle$ be a TDG. The TDG-c estimates of $\mathcal{G}$ 's task nodes are given by:

$$h_{\mathcal{T}}(v_t) := egin{cases} cost(v_t) & ext{if } v_t ext{ primitive} \ \min_{(v_t,v_m)\in \mathcal{E}_{\mathcal{T}
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Cost-sensitive	TDG Heuristic, TDG-c			
TDG-c C	Computation			

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For methods nodes  $v_m = (T, \prec, VC, \alpha)$ :

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Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○○○○○○○○○○○○○○○	Compilation Technique	
Cost-sensitive	TDG Heuristic, TDG-c			
TDG-c C	Computation			

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Heuristic value for  $tn = (T, \prec, VC, \alpha)$ :

decomposition-based: 
$$h_{TDG-c}(tn) := \sum_{\substack{t(\bar{\tau}) \in T \\ t(\bar{\tau}) \text{ abstract}}} \left( \min_{v_t \in Ground_{VC}(t(\bar{\tau}))} h_T(v_t) \right)$$



Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○○○●○○○○○○○○○○	Compilation Technique	
Cost-sensitive	TDG Heuristic, TDG-c			
TDG-c C	Computation			

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progression-based: 
$$h_{TDG-c}(tn) := \sum_{t(\bar{\tau})\in T} \left( \min_{v_t \in Ground_{VC}(t(\bar{\tau}))} h_T(v_t) \right)$$

1

1

Introduction	Task Decomposition Graph		TDG-c & TDG-m ○○OOO●○○○○○○○○○	Compilation Technique	
Cost-sensitive	ooooooooooooooooooooooooooooooooooooo				
Notos					

$$\sum_{\substack{t(\bar{\tau})\in T\\t(\bar{\tau}) \text{ abstract}}} \left(\min_{v_t \in Ground_{VC}(t(\bar{\tau}))} h_T(v_t)\right) \text{ becomes } \sum_{\substack{t(\bar{c})\in T\\t(\bar{c}) \text{ abstract}}} h_T(t(\bar{c})).$$



	Task Decomposition Graph	TDG-c & TDG-m ○○○○○●○○○○○○○○○	Compilation Technique	
Cost-sensitive	TDG Heuristic, TDG-c			
ALC: N				

The heuristic formulae were given for *lifted* planning and can be simplified for ground planning:

$$\sum_{\substack{t(\bar{\tau})\in T\\t(\bar{\tau}) \text{ abstract}}} \left( \min_{v_t \in Ground_{VC}(t(\bar{\tau}))} h_T(v_t) \right) \text{ becomes } \sum_{\substack{t(\bar{c})\in T\\t(\bar{c}) \text{ abstract}}} h_T(t(\bar{c})).$$

Why do we have two different formulae for progression- and decomposition-based search?



	Task Decomposition Graph	TDG-c & TDG-m ○○○○○●○○○○○○○○○○	Compilation Technique	
Cost-sensitive	TDG Heuristic, TDG-c			

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- Why do we have two different formulae for progression- and decomposition-based search?
  - Because the costs of search nodes are usually computed differently.



Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○ <b>○○○</b> ●○○○○○○○○○	Compilation Technique	
Cost-sensitive	TDG Heuristic, TDG-c			

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Introduction	Task Decomposition Graph		TDG-c & TDG-m ○○○○○●○○○○○○○○○○	Compilation Technique				
Cost-sensitive TDG Heuristic, TDG-c								

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Introduction	Task Decomposition Graph		TDG-c & TDG-m ○○○○○●○○○○○○○○○○	Compilation Technique				
Cost-sensitive TDG Heuristic, TDG-c								

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  - Because the costs of search nodes are usually computed differently.
  - In decomposition-based search, a search node's costs is given by its primitive tasks.
  - In progression-based search, those costs are usually incorporated after an action has been progressed away.
  - → Using only the abstract tasks for the heuristic in decomposition-based planning thus prevents taking those primitive costs into account twice.





Just estimating the final solution costs says little about the effort finding it. One can easily construct examples, where expensive solutions can be found easily (with only few decompositions), whereas cheap solutions need more search effort.



Introduction	Task Decomposition Graph		TDG-c & TDG-m ○○○○○○●○○○○○○○○○	Compilation Technique			
Modification-sensitive TDG Heuristic, TDG-m							
Motivatio	n						

- Just estimating the final solution costs says little about the effort finding it. One can easily construct examples, where expensive solutions can be found easily (with only few decompositions), whereas cheap solutions need more search effort.
- We thus exploit the TDG to estimate how many *modifications* we require for certain tasks.



Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○○○○○●○○○○○○○○○	Compilation Technique	
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- We thus exploit the TDG to estimate how many *modifications* we require for certain tasks.
- The resulting heuristic will be not be admissible, but admissible in the number of required modifications (trivial). This means that any solution returned by A\* will have the property that no other solution can be created with fewer modifications. (This is not something we aim for, it's just a property we get.)

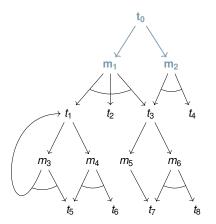


TDG-c & TDG-m 

Modification-sensitive TDG Heuristic, TDG-m

Illustration of TDG-m Computation - For Decomposition-based Search

**Example:** 



$$h_T(t_0) = 1 + \min\{h_M(m_1), h_M(m_2)\}$$

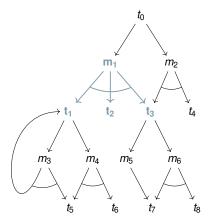


Introduction Task Decomposition Graph Landmarks TDG-c & TDG-m Compilation Technique Summary

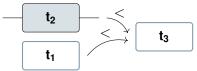
Modification-sensitive TDG Heuristic, TDG-m

Illustration of TDG-m Computation – For Decomposition-based Search

# Example:



Method  $m_1 = (t_0, t_n)$  with task network *tn*:



$$h_M(m_1) = \sum_{t_i \in \{t_1, t_2, t_3\}} h_T(t_i)$$

(Also subtract |CL| in case we have a partial plan containing causal links.)

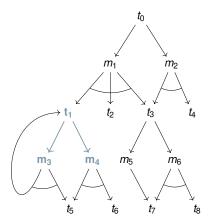


TDG-c & TDG-m 

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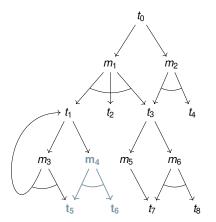


Introduction Task Decomposition Graph Landmarks TDG-c & TDG-m Compilation Technique Summary

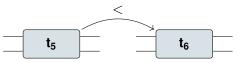
Modification-sensitive TDG Heuristic, TDG-m

Illustration of TDG-m Computation - For Decomposition-based Search

Example:



Method  $m_4 = (t_1, t_n)$  with task network *tn*:



$$egin{aligned} h_M(m_4) &= h_T(t_5) + h_T(t_6) \ &= |pre(t_5)| + |pre(t_6)| \ &= 2+2 = 4 \end{aligned}$$

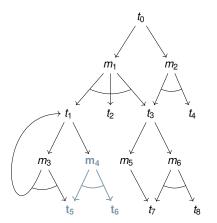


Introduction Task Decomposition Graph Landmarks TDG-c & TDG-m Compilation Technique Summary

Modification-sensitive TDG Heuristic, TDG-m

Illustration of TDG-m Computation – For Decomposition-based Search

Example:



Method  $m_4 = (t_1, P)$  with partial plan P:



 $h_{M}(m_{4}) = h_{T}(t_{5}) + h_{T}(t_{6}) - |CL|$ = |pre(t\_{5})| + |pre(t\_{6})| - 2 = 2 + 2 - 2 = 2



	Task Decomposition Graph	TDG-c & TDG-m ○○○○○○○○○○○○○○○○	Compilation Technique	
Modification-se	ensitive TDG Heuristic, TDG-m			
TDO (	0 1 1			

Let  $\mathcal{G} = \langle V_T, V_M, E_{T \to M}, E_{M \to T} \rangle$  be a TDG. The TDG-c estimates of  $\mathcal{G}$ 's task nodes are given by:

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	Task Decomposition Graph	TDG-c & TDG-m ○○○○○○○○○○○○○○○○	Compilation Technique	
Modification-se	ensitive TDG Heuristic, TDG-m			
TDO	0			

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For methods nodes  $v_m = (T, \prec, VC, CL, \alpha)$ :

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Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○○○○○○○○○○○○○○○○	Compilation Technique	
Modification-se	ensitive TDG Heuristic, TDG-m			
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Heuristic value for partial plan  $P = (T, \prec, VC, CL, \alpha)$ : decomposition-based:  $h_{TDG-m}(P) := \sum_{t(\bar{\tau})\in T} \left( \min_{v_t \in Ground_{VC}(t(\bar{\tau}))} h_T(v_t) \right) - |CL|$ 



Introduction	Task Decomposition Graph		TDG-c & TDG-m ○○○○○○○○○○○○○○○○	Compilation Technique			
Modification-sensitive TDG Heuristic, TDG-m							
TDO	0						

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Heuristic value for partial plan  $P = (T, \prec, VC, CL, \alpha)$ : decomposition-based:  $h_{TDG-m}(P) := \sum_{t(\bar{\tau})\in T} \left( \min_{v_t \in Ground_{vC}(t(\bar{\tau}))} h_T(v_t) \right) - |CL|$ progression-based: Ignore links and use 1 instead of  $|pre(v_t)|$ .

Chapter: Heuristics for (Hierarchical) Planning Problems by Dr. Pascal Bercher

Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○○○○○○○●○○○○○○	Compilation Technique	
Properties of T	DG Heuristics			
Propertie	s			

Note that this heuristic does not compute executable plans. So what *does* it compute?



	Task Decomposition Graph	TDG-c & TDG-m ○○○○○○○○●○○○○○○	Compilation Technique	
Properties of T	DG Heuristics			
Propertie	es			

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Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○○○○○○○●○○○○○○	Compilation Technique	
Properties of 7	TDG Heuristics			
Propertie	es			

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Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○○○○○○○●○○○○○○	Compilation Technique	
Properties of T	DG Heuristics			
Propertie	es			

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- However, the costs (or the effort) of these tasks is not reflected in the heuristic value!
- To illustrate what this means: What heuristic do we get if the only abstract task can be decomposed into an empty task network (which will not work as solution due to a goal description or subsequent primitive tasks).



Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○○○○○○○●○○○○○○	Compilation Technique	
Properties of 7	TDG Heuristics			
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- However, the costs (or the effort) of these tasks is not reflected in the heuristic value!
- To illustrate what this means: What heuristic do we get if the only abstract task can be decomposed into an empty task network (which will not work as solution due to a goal description or subsequent primitive tasks).
- However, the heuristic can still come up with even exponentially large heuristic values. (This is true although every task occurs just once in the TDG. Why?)





Recall from the beginning (construction of the TDG) that the TDG can be recomputed as soon as any of its tasks is identified as unreachable.



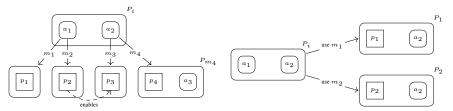


- Recall from the beginning (construction of the TDG) that the TDG can be recomputed as soon as any of its tasks is identified as unreachable.
- So far, the TDG is only computed *once*. However, different search nodes might have different reachable action sets!





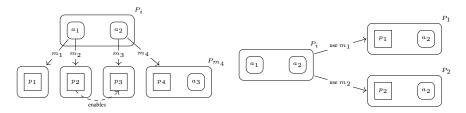
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Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○○○○○○○○○○○○○○○○○	Compilation Technique	
TDG Recompu	itation			

### Example

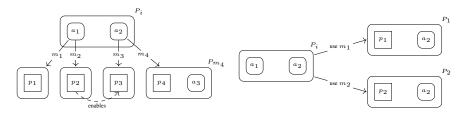


• Let  $c(p_3) = i$  and  $h_M(P_{m_4}) = h_T(p_4) + h_T(a_3) = j > i$ .



Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○○○○○○○○○●○○○○	Compilation Technique	
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#### Example

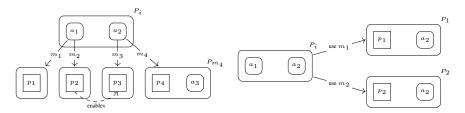


- Let  $c(p_3) = i$  and  $h_M(P_{m_4}) = h_T(p_4) + h_T(a_3) = j > i$ .
- Then, we get  $h_T(a_2) = i$ . Let us now consider the heuristic values for  $P_1$  and  $P_2$  resulting from decomposing  $a_1$  using  $m_1$  or  $m_2$ , respectively.



	Task Decomposition Graph		TDG-c & TDG-m ○○○○○○○○○○●○○○○	Compilation Technique			
TDG Recomputation							

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- Without recomputation, we get  $h(P_1) = h(P_2) = i$ . With recomputation, we get  $h(P_1) = j$  and  $h(P_2) = i$ , so we get improved heuristic accuracy due to updated reachability information in the TDG.





When to Recompute? - In Decomposition-based Planning

Let *P* be a partial plan, *mod* a modification, and *P'* the partial plan resulting from applying *mod* to *P*.





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- Why? If there is just one method, the reachable action set cannot possibly change.
- There are other cases, which are not yet handled, though. E.g., causal links might also limit the available actions.



Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○○○○○○○○○○○○○	Compilation Technique	
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When to Recompute? – In Progression-based Planning

Here, we have just two modifications: method application and progression (i.e., action application).





When to Recompute? – In Progression-based Planning

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#### When to Recompute? – In Progression-based Planning

- Here, we have just two modifications: method application and progression (i.e., action application).
- With methods, we have the same situation as in decomposition-based planning (since also here, decompositions restrict the reachable actions).
- Also, each progression allows a recomputation.





If mod is not a decomposition (i.e., an insertion of a causal link, an ordering, or a variable constraint), we get:  $h_{TDG-c}(P) = h_{TDG-c}(P')$ 





# What if we don't Recompute TDG-c?

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- If mod is a method m = (t, tn<sub>m</sub>) (without alternatives), we can set:

$$h_{TDG-c}(P') = h_{TDG-c}(P) - \sum_{\substack{t(\bar{c}) \in T_m \\ t(\bar{c}) \text{ primitive}}} c(t(\bar{c}))$$





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These equations are specific to decomposition-based planning. The latter is required because the method's primitive tasks were accounted by the heuristic, but are now covered by the cost value of the search node.



Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○○○○○○○○○○○○○	Compilation Technique	
TDG Recomp	utation			

What if we don't Recompute TDG-m?

If *mod* is an ordering or variable insertion, we get:  $h_{TDG-m}(P) = h_{TDG-m}(P')$ 



	Task Decomposition Graph	TDG-c & TDG-m ○○○○○○○○○○○○○○	Compilation Technique	
TDG Recompu	utation			

# What if we don't Recompute TDG-m?

- If *mod* is an ordering or variable insertion, we get:  $h_{TDG-m}(P) = h_{TDG-m}(P')$
- If mod is a causal link insertion or a decomposition (without alternatives), we get:  $h_{TDG-m}(P) = h_{TDG-m}(P') 1$

Introduction	Task Decomposition Graph	TDG-c & TDG-m ○○○○○○○○○○○○○○	Compilation Technique	
TDG Recompu	itation			

# What if we don't Recompute TDG-m?

- If *mod* is an ordering or variable insertion, we get:  $h_{TDG-m}(P) = h_{TDG-m}(P')$
- If mod is a causal link insertion or a decomposition (without alternatives), we get:
   h<sub>TDG-m</sub>(P) = h<sub>TDG-m</sub>(P') 1
- No difference between decomposition- and progression-based planning.



Introduction	Task Decomposition Graph	TDG-c & TDG-m 000000000000000000	Compilation Technique ●000000000	
Motivatio	าท			

We would like to exploit existing classical planning heuristics in HTN planning.



	Task Decomposition Graph	TDG-c & TDG-m 00000000000000000000	Compilation Technique ●○○○○○○○○○	
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	Task Decomposition Graph	TDG-c & TDG-m 000000000000000000	Compilation Technique	
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Motivation

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- Issues:
  - More expressive formalism (cf. lecture on expressiveness and next lecture on computational complexity). In particular: Hierarchical problems are undecidable, so they cannot be translated into classical problems.



Introduction	Task Decomposition Graph	TDG-c & TDG-m 0000000000000000000	Compilation Technique •000000000	

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Introduction	Task Decomposition Graph	TDG-c & TDG-m 00000000000000000	Compilation Technique •00000000	

## Motivation

- We would like to exploit existing classical planning heuristics in HTN planning.
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  - In general, HTN problems do not have a goal description, so what's the classical problem's goal?
  - In classical planning, all actions can be applied, in hierarchical planning only those reachable from the initial task (network).



Introduction	Task Decomposition Graph	TDG-c & TDG-m 00000000000000000	Compilation Technique	

Similarities and differences to TDG heuristic:

Both heuristics can estimate plan cost or modification effort.



Introduction	Task Decomposition Graph	TDG-c & TDG-m 00000000000000000	Compilation Technique	

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Task Decomposition Graph	TDG-c & TDG-m 00000000000000000	Compilation Technique	

Similarities and differences to TDG heuristic:

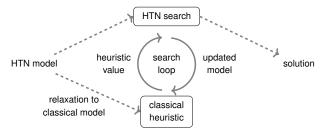
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- They both, somehow, incorporate the TDG.

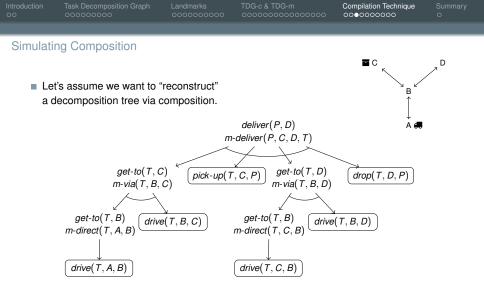


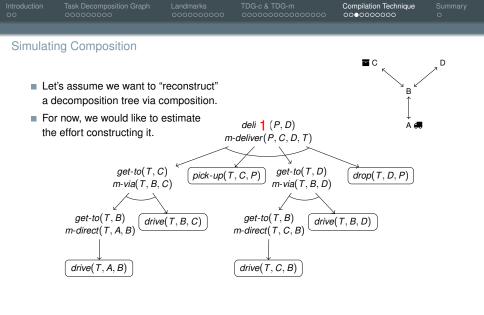


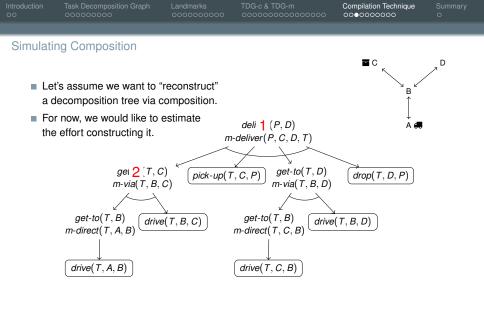
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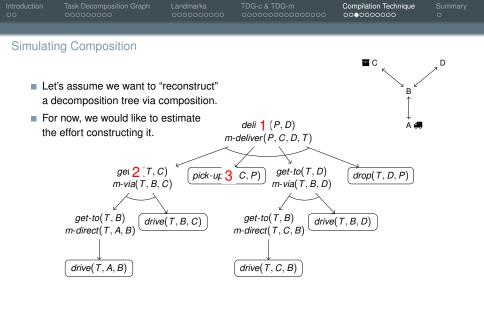
- Both heuristics can estimate plan cost or modification effort.
- Both heuristics work for progression- and decomposition-based planning.
- They both, somehow, incorporate the TDG.
- It is not a preprocessing heuristic: Its "heuristics model" gets adapted for every search node. In that way, it corresponds the previous heuristics with enabled recomputation.

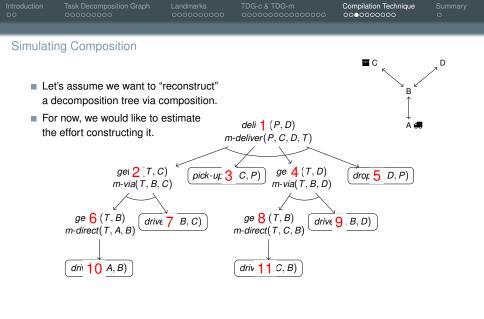


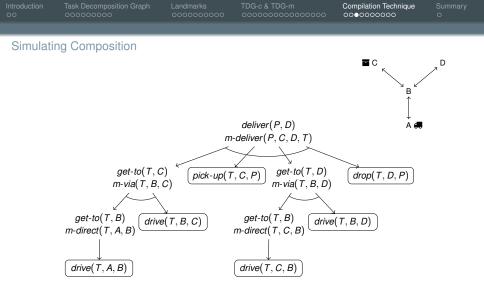


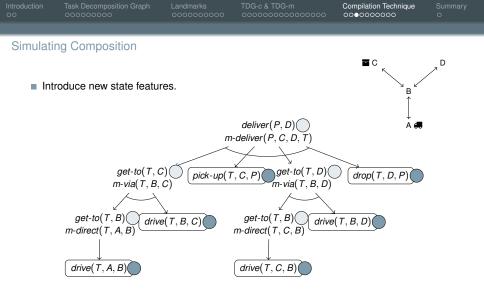


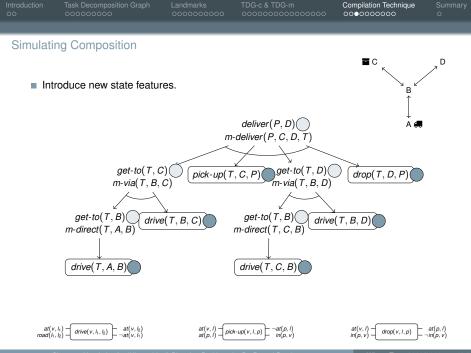




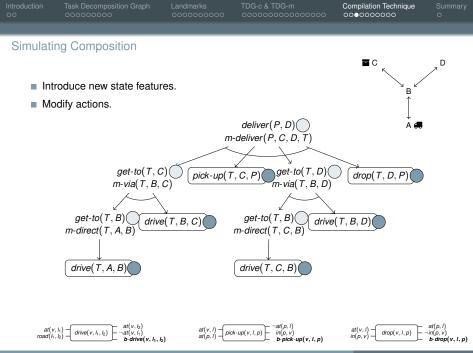




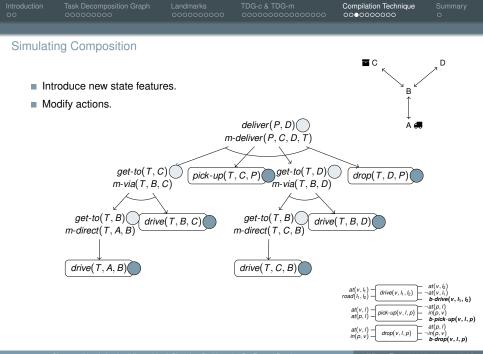


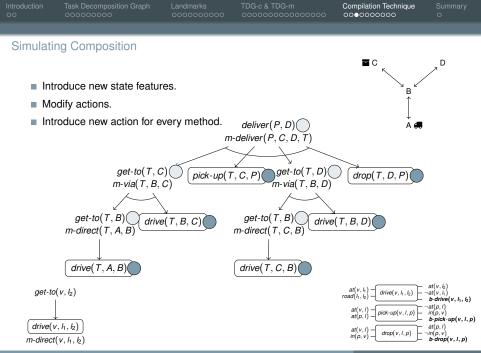


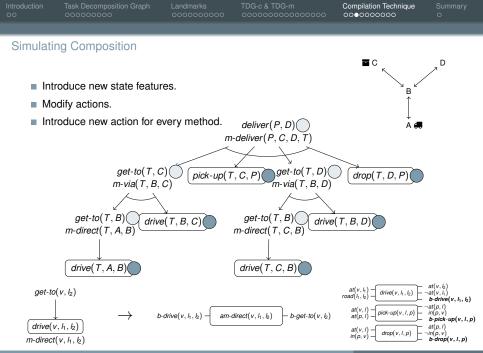
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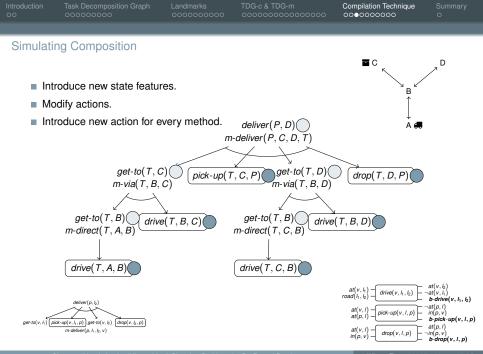
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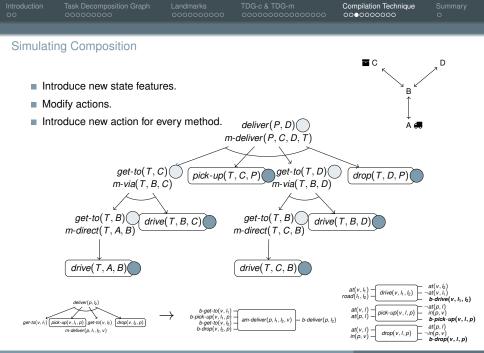




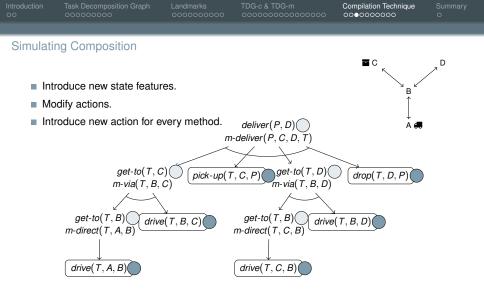


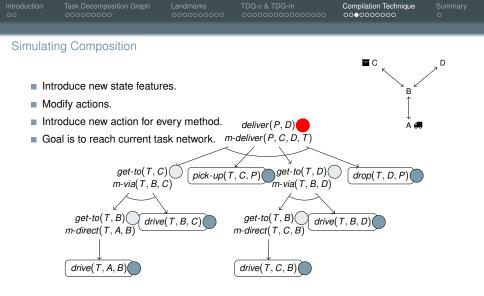
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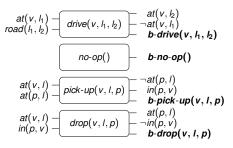




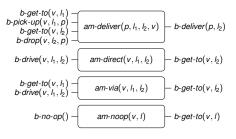
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## Simulating Composition – Resulting Model

# Altered action encodings:



# New actions encoding methods:





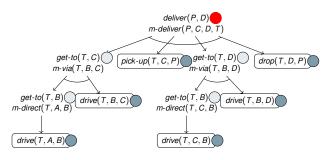
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Summary

# Unrelaxed Planning in the Transformed Model





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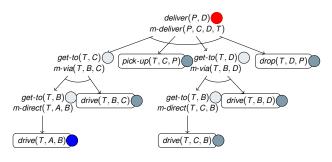
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Summary

## Unrelaxed Planning in the Transformed Model





 $\begin{array}{c} \{at(T,A), \\ at(P,C)\} \end{array} \qquad \begin{array}{c} \{at(T,B), \\ drive(T,A,B) \\ b \cdot drive(T,A,B)\} \end{array}$ 

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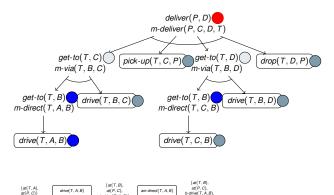
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## Unrelaxed Planning in the Transformed Model

b-drive(T, A, B)}





b-get-to(T,B)

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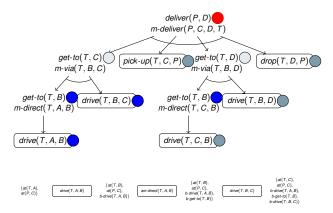
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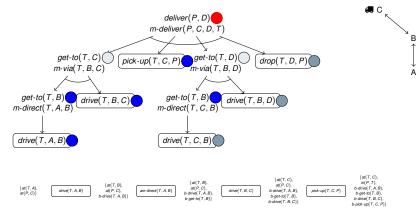




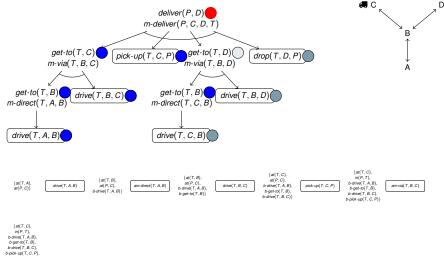
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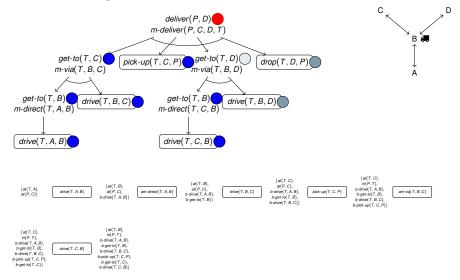


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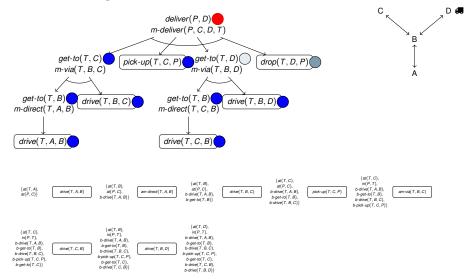


b-get-to(T,C)}

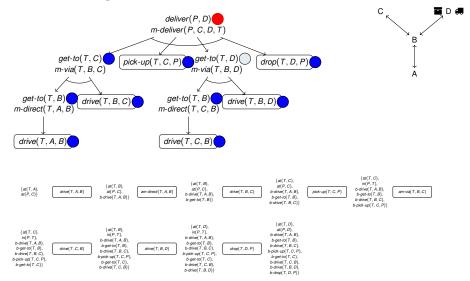
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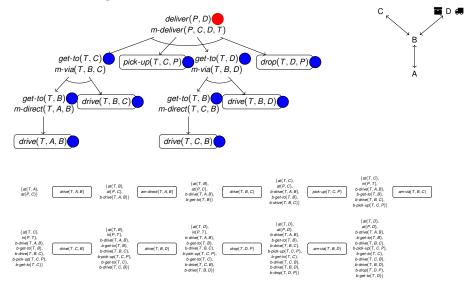
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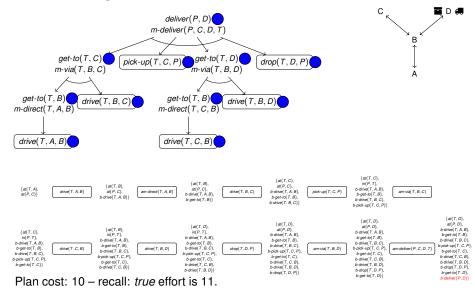
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Introduction

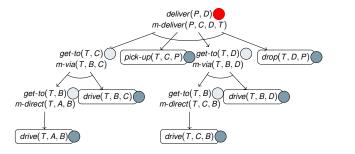
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Summary

### Relaxed Planning in the Transformed Model (Heuristic Computation)





 $\{at(T, A), at(P, C)\}$ 

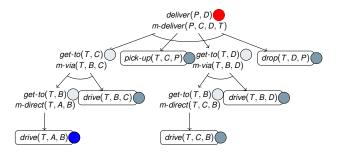


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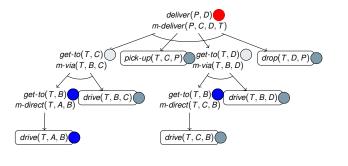


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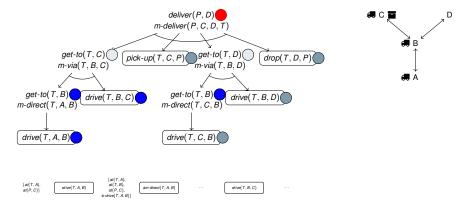


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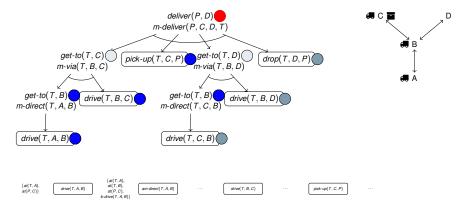


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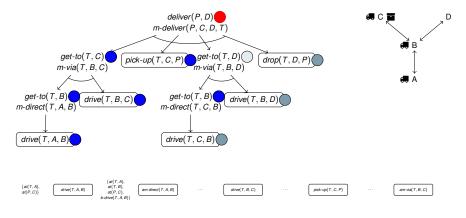


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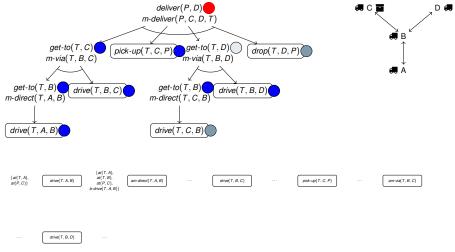


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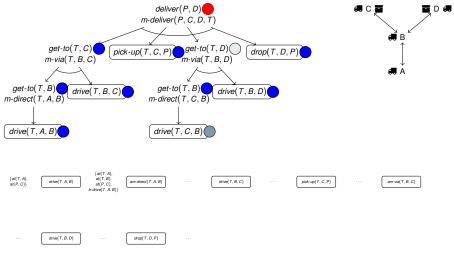
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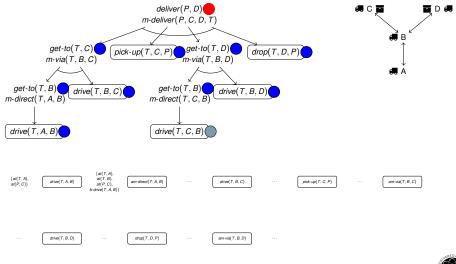


sk Decomposition Graph

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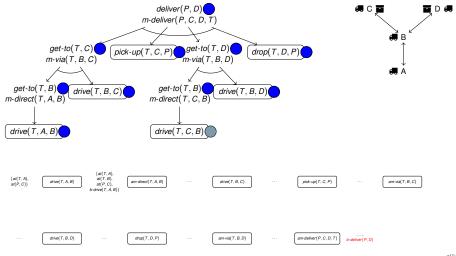
sk Decomposition Graph

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DG-c & TDG-m 0000000000000000 Compilation Technique

Summary

# Relaxed Planning in the Transformed Model (Heuristic Computation)



Plan cost: 9 - recall: true effort is 11 and unrelaxed plan cost is 10.



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Introduction	Task Decomposition Graph	TDG-c & TDG-m 000000000000000000	Compilation Technique	
General	Characteristics			

# Technique simulates task composition,



Introduction	Task Decomposition Graph	TDG-c & TDG-m 000000000000000000000000000000000000	Compilation Technique	

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- ✓ incorporates hierarchical reachability information,
- $\checkmark$  combines it with information on state-based executability, and
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    - HTN ordering relations are relaxed.
  - Heuristic function may only insert tasks that lie within the decomposition hierarchy (not given here).



Introduction	Task Decomposition Graph	TDG-c & TDG-m 0000000000000000	Compilation Technique	

The size of the new model is *linear* in the input HTN domain.



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  - Admissibility (only if costs are chosen as above).



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Discussio	on			

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Introduction	Task Decomposition Graph		TDG-c & TDG-m 000000000000000000	Compilation Technique 000000000●		
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  - Many classical heuristics compute heuristic values that are polynomial in the the input. The TDG heuristics can come up with exponential heuristic values.



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Summar	У			

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- Landmarks tasks that occur on any sequence from the initial task network to a solution – can be used as basis for heuristics.
- The TDG heuristics compute admissible estimates, but take task insertion into account to only a limited extent.
- We can also exploit *classical* heuristics for hierarchical planning by a relaxing problem transformation.

