COMP3630 / COMP6363

week 2: **Properties of Regular Languages**

This Lecture Covers Chapter 4 of HMU: Properties of Regular Languages

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The Australian National University

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Content of this Chapter

- Pumping Lemma for regular languages
- Some more properties of regular languages
- Decision properties of regular languages
- Equivalence and minimization of automata

Additional Reading: Chapter 4 of HMU.

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Lemma 4.1.1 (Pumping Lemma for Regular Languages)

Let L be a regular language. There there exists an $n \in \mathbb{N}, n \geq 1$ (depending on L) such that for any string $w \in L$ with $|w| \geq n$, there exist strings x, y, z such that:
(1) $w = xyz$
(2) $|xy| \leq n$
(3) $|y| > 0$
(4) $xy^i z \in L$ for $i \in \mathbb{N} \cup \{0\}$

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- Then, there must be a sequence of transitions that move A from q_0 to some final state upon reading w.

$$\underbrace{q_0 = q_{i_0} \longrightarrow q_{i_1} \longrightarrow q_{i_2} \longrightarrow \cdots \longrightarrow q_{i_n}}_{n \text{ symbols and } n+1 \text{ states}} \longrightarrow \cdots \longrightarrow q_{i_k} \in F$$

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- **SOME** state must be visited (at least) twice. Let $q_{i_a} = q_{i_b}$ for $i_0 \leq i_a < i_b \leq i_n$.

$$q_0 = q_{i_0} \overbrace{\xrightarrow{s_1} \cdots \xrightarrow{s_{i_a}}}^{x \text{ is read}} {\color{blue}q_{i_a}} \overbrace{\xrightarrow{s_{i_a}+1} \cdots \xrightarrow{s_{i_b}}}^{y \text{ is read}} {\color{blue}q_{i_b}} \overbrace{\xrightarrow{s_{i_b}+1} \cdots \xrightarrow{s_{i_n}} q_{i_n} \xrightarrow{s_{i_n}+1} \cdots \xrightarrow{s_{i_k}}}^{z \text{ is read}} q_{i_k} \in F$$

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- (4) holds since the path for xy^iz is derived from the above either by deleting the subpath between q_{i_a} and q_{i_b} or by repeating it. All such paths end in $q_{i_k} \in F$.

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  - let's apply this to the string $w = 0^k 1^k \in L$:

$$\underbrace{0\ldots0}_{x}\underbrace{0\ldots0}_{y}\underbrace{0\ldots0\,1\ldots1}_{z}\quad\text{with}\quad\overbrace{0\ldots0\,0\ldots0\,0\ldots0}^{k}\overbrace{1\ldots1}^{k}$$

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- $L = \{w \in \{0,1\}^* : |w| \text{ is a prime}\}$ is **not** regular.
- $L = \{ww^R : w \in \{0,1\}^*\}$ is **not** regular. [$w^R = w$ read from right to left].

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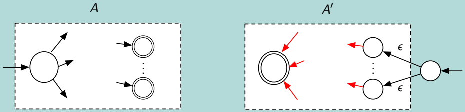
Proof of Theorem 4.2.1

- Let langauge L be accepted by DFA A.



- Let A' be the DFA obtained by: (a) Reversing each arrow in A; (b) swapping final and initial states; and (c) introduce ϵ-transitions to make initial state (of A') unique.
- Then L^R is accepted by A'.

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Expressions Languages

$$\begin{tikzcd} E \arrow[r] \arrow[d, "h"'] & L(E) \arrow[d, "h"] & \text{(language over } \Sigma_1) \\ h(E) \arrow[r] & h(L(E)) = L(h(E)) & \text{(language over } \Sigma_2) \end{tikzcd}$$

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- Hence, B accepts $h^{-1}(L)$.

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  - Given two DFAs A and B, is $L(A) \subseteq L(B)$?
- We will look at the above 5 questions assuming that regular languages are defined by DFAs. (If the language is specified by an expression, we can convert it to a DFA!)

Decision Properties (Emptiness, Membership, Infiniteness)

- **Emptiness:** If one is given a DFA with n states that accepts L, we can find all the states reachable from the initial state in $O(n^2)$ time. If no final state is reachable, L must be empty.

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- **Infiniteness:** We can reduce the problem of infiniteness to finding a cycle in the directed graph (a.k.a. transition diagram) of the DFA.
  - First, delete any node unreachable from the initial node ($O(n^2)$ complexity).
  - Next, delete nodes that cannot reach any final node ($O(n^3)$ complexity).
  - Use depth-first search (DFS) to find a cycle in the remaining graph ($O(n^2)$ complexity).

Decision Properties (Emptiness, Membership, Infiniteness)

- **Emptiness:** If one is given a DFA with n states that accepts L, we can find all the states reachable from the initial state in $O(n^2)$ time. If no final state is reachable, L must be empty.
- **Membership:** If one is given a DFA with n states that accepts L, given string w, we can simply identify the transitions corresponding to w one symbol at a time. If the last state is an accepting state, then w must be in the language. This takes no more than $O(|w|)$ time steps.
- **Infiniteness:** We can reduce the problem of infiniteness to finding a cycle in the directed graph (a.k.a. transition diagram) of the DFA.
  - First, delete any node unreachable from the initial node ($O(n^2)$ complexity).
  - Next, delete nodes that cannot reach any final node ($O(n^3)$ complexity).
  - Use depth-first search (DFS) to find a cycle in the remaining graph ($O(n^2)$ complexity).
- **Q:** How do runtimes change if we have an NFA?

Decision Properties (Equivalence)

- **Equivalence:** Given $A = (Q_A, \Sigma, \delta_A, q_{A0}, F_A)$ and $B = (Q_B, \Sigma, \delta_B, q_{B0}, F_B)$, how do we ascertain if $L(A) = L(B)$?

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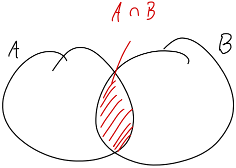
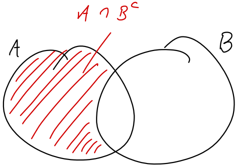
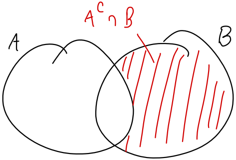
$$L(A) = L(B) \Leftrightarrow \begin{array}{l} L(A) \cap L(B)^c = \emptyset \\ L(A)^c \cap L(B) = \emptyset \end{array}$$

Why is this true?







Thus:

$A \cap B = A = B$ iff

$A \cap B^c = A^c \cap B = \emptyset$

Decision Properties (Equivalence)

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Run A and B in parallel. (Not their complement-versions!)

  - $L(A) \cap L(B)^c$: Accept if resp. paths end in F_A and F_B^c.
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- Use **product** DFA: Construct $C = (Q_C, \Sigma, \delta_C, q_{C0}, F_C)$ defined by

$$\begin{aligned} Q_C &= Q_A \times Q_B && \text{[Cartesian Product]} \\ q_{C0} &= (q_{A0}, q_{B0}) && \\ \delta_C((q, q'), s) &= (\delta_A(q, s), \delta_B(q', s)) && \text{[Both DFAs are simulated in parallel]} \\ F_C &= (F_A \times F_B^c) \cup (F_A^c \times F_B) && \text{[accept strings in exactly one of } L(A) \text{ or } L(B)\text{]} \end{aligned}$$

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$$L(A) = L(B) \Leftrightarrow L(C) = \emptyset$$

(and you know how to check for an empty language of a DFA!)

Decision Properties (Inclusion)

- **Inclusion:** Given $A = (Q_A, \Sigma, \delta_A, q_{A0}, F_A)$ and $B = (Q_B, \Sigma, \delta_B, q_{B0}, F_B)$, how do we ascertain if $L(A) \subseteq L(B)$?

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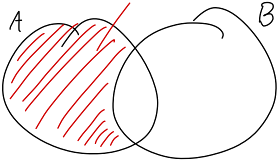
Why is this true?

$A \cap B^c = \emptyset$



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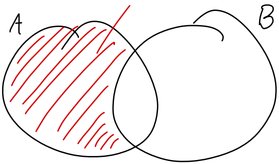
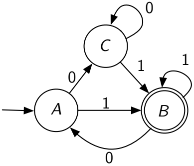


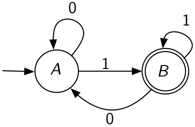
Thus: $A \subseteq B$

What now? Everything as before! Only now $F_C = F_A \times F_B^c$.

DFA State Minmimization

- Given two DFAs, we know how to test if they accept the same language.
- Is there a unique minimal DFA for a given regular language?
- Given a DFA, can we **reduce** the number of states without altering the language it accepts?





Clearly, the two DFAs accept the same language and state C is unnecessary.

- How do we (identify and) remove 'unnecessary' states without altering the underlying language?

DFA State Minimimization

- State minimization requires a notion of **equivalence** or **distinguishability** of states.

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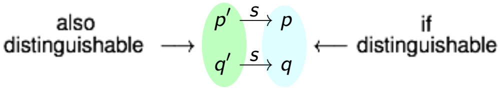
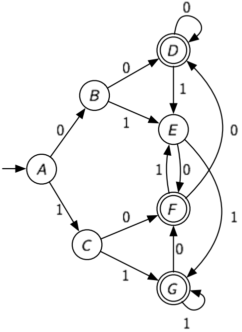
- **Table Filling** Algorithm identifies equivalent and distinguishable pairs of states.
  - Any final state is **distinguishable** from a non-final state (and vice versa)
  - If (a) p and q are distinguishable; and there exist states p', q', and symbol s such that (b) $\hat{\delta}(p', s) = p$ and (c) $\hat{\delta}(q', s) = q$, then p' and q' are also distinguishable



Identifying pairs of (In)distinguishable States: An Example



	G	F	E	D	C	B	A
A							
B							
C							
D							
E							
F							
G							

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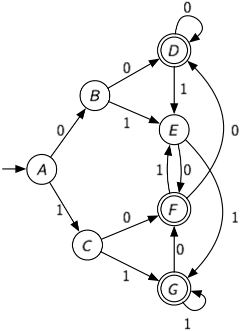

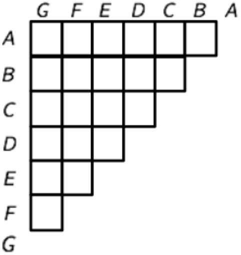

- Fill in $\times$ whenever one component of pair is final, and other is not.

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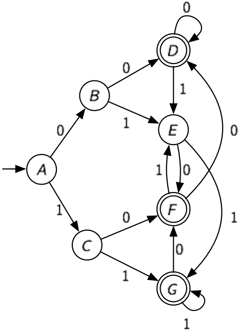

	G	F	E	D	C	B	A
A							
B							
C							
D							
E							
F							
G							

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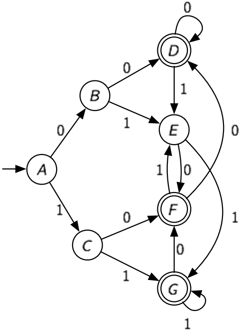

	G	F	E	D	C	B	A
A							
B							
C							
D							
E							
F							
G							

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- Fill in × if 0 moves the pair of states to a distinguishable pair

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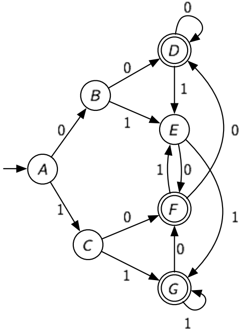

	G	F	E	D	C	B	A
A							
B							
C							
D							
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(This slide is added to the handout so you can try it yourself!)

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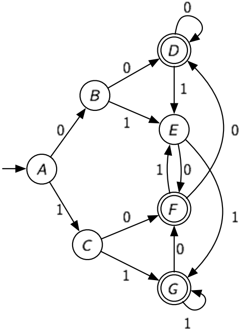

	G	F	E	D	C	B	A
A	×	×	×	×	×	×	
B	×	×	×	×	×		
C	×	×		×			
D	×		×				
E	×	×					
F	×						
G							

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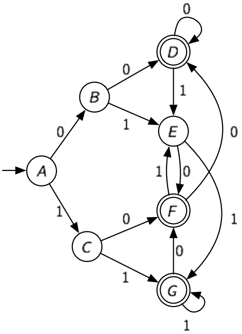

	G	F	E	D	C	B	A
A	×	×	×	×	×	×	
B	×	×	×	×	×		
C	×	×		×			
D	×		×				
E	×	×					
F	×						
G							

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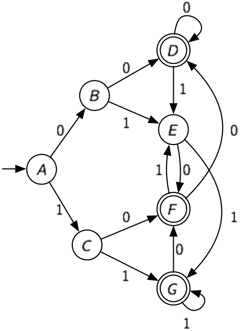
Theorem 4.4.1

Any two states without a $\times$ sign are equivalent.

Identifying pairs of (In)distinguishable States: An Example



	G	F	E	D	C	B	A
A	×	×	×	×	×	×	
B	×	×	×	×	×		
C	×	×		×			
D	×		×				
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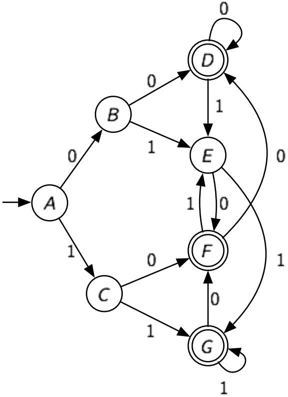
Theorem 4.4.1

Any two states without a × sign are equivalent.

- Proof idea: If two states are distinguishable, the algorithm **will** fill a × eventually.

Table-filling Algorithm



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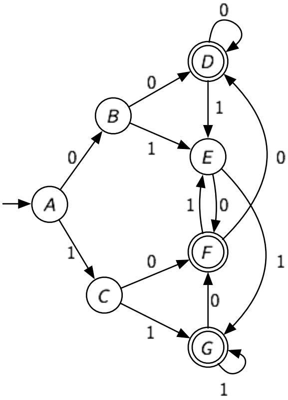

- Delete states not reachable from start states

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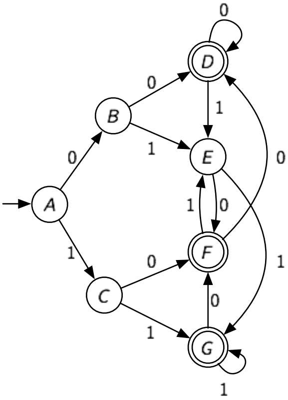
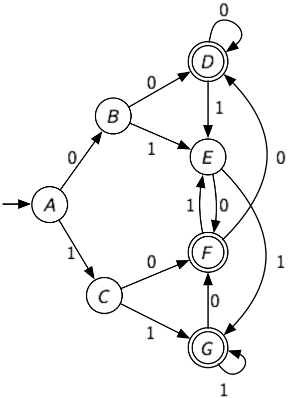


- Delete states not reachable from start states
- Delete any non-starting state that cannot reach any final state

Table-filling Algorithm



- Delete states not reachable from start states
- Delete any non-starting state that cannot reach any final state
- Find distinguishable and equivalent pairs of states

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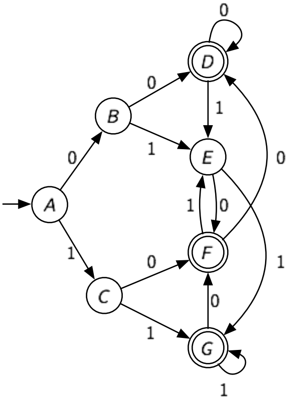

- Delete states not reachable from start states
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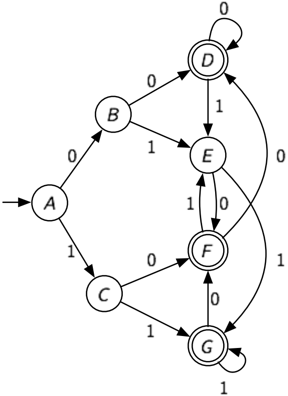
- Find distinguishable and equivalent pairs of states
- Find equivalence classes of indistinguishable states. In this example: $\{A\}, \{B\}, \{C, E\}, \{D, F\}, \{G\}$

	G	F	E	D	C	B	A
A	×	×	×	×	×	×	
B	×	×	×	×	×		
C	×	×		×			
D	×		×				
E	×	×					
F	×						
G							

Color-blind?
red: A/B, blue: A/E, A/C, B/E, B/C, D/G, F/G

Table-filling Algorithm



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- Collapse each equivalence class of states to a state

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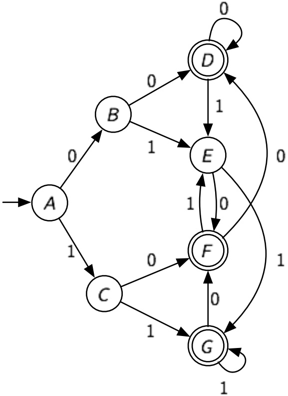

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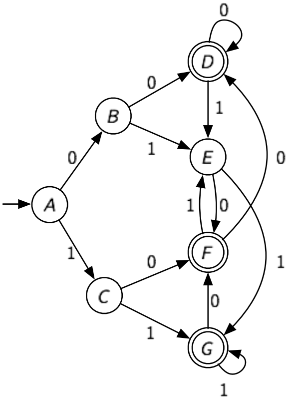
- Delete parallel transitions with same label.

Table-filling Algorithm



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Remark: The resultant transition diagram will be a DFA.

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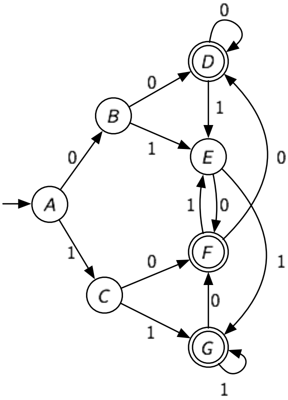

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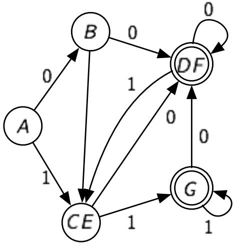

Table-filling: Other Uses

- Test equivalence of languages accepted by 2 DFAs.

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    - Rename states in Q_B so that Q_A and Q_B are disjoint.

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      (Ignore the fact that there are 2 start states)

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    - q_{A0} and q_{B0} are indistinguishable $\Leftrightarrow L(A) = L(B)$.

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      (Ignore the fact that there are 2 start states)
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    - q_{A0} and q_{B0} are indistinguishable $\Leftrightarrow L(A) = L(B)$.

  [Why?] If w distinguishes q_{A0} from q_{B0} then w cannot be in both $L(A)$ and $L(B)$

Table-filling: Other Uses

- Test equivalence of languages accepted by 2 DFAs.
  - Given $A = (Q_A, \Sigma, \delta_A, q_{A0}, F_A)$ and $B = (Q_B, \Sigma, \delta_B, q_{B0}, F_B)$:
    - Rename states in Q_B so that Q_A and Q_B are disjoint.
    - View A and B together as one DFA
      (Ignore the fact that there are 2 start states)
    - Run table-filling on $Q_A \cup Q_B$.
    - q_{A0} and q_{B0} are indistinguishable $\Leftrightarrow L(A) = L(B)$.

  [Why?] If w distinguishes q_{A0} from q_{B0} then w cannot be in both $L(A)$ and $L(B)$
- Suppose a DFA A cannot be minimized further by table-filling. Then, A has the least number of states among all DFAs that accept $L(A)$