

On Delete Relaxation in Partial-Order Causal-Link Planning

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Outline

(**POCL** Planning == **P**artial-**O**rder **C**ausal-**L**ink Planning)

- 1 POCL Planning
- 2 Delete Relaxation in POCL Planning
- 3 Complexity Results
 - NP Membership
 - NP Hardness
- 4 A new Heuristic for POCL Planning
 - Sample FF: Heuristic Function
 - Evaluation
- 5 Summary and Outlook



A planning domain is a tuple $\mathcal{D} = \langle \mathcal{V}, \mathcal{A} \rangle$ with:

- \mathcal{V} is a finite set of state variables, $s \in 2^{\mathcal{V}}$ being a state,
- \mathcal{A} is a finite set of actions,
an action $a := \langle pre, add, del \rangle \in \mathcal{A}$ consists of:
 - $pre \subseteq \mathcal{V}$, the precondition of a ,
 - $add \subseteq \mathcal{V}$, the add list of a ,
 - $del \subseteq \mathcal{V}$, the delete list of a



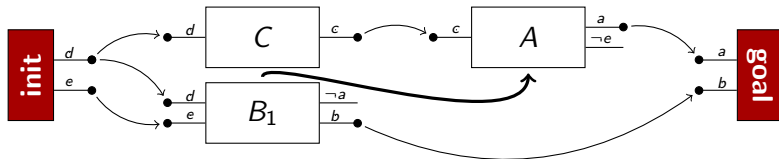
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- \mathcal{A} is a finite set of actions,
an action $a := \langle pre, add, del \rangle \in \mathcal{A}$ is applicable:
 - in a state $s \in 2^{\mathcal{V}}$ iff $pre \subseteq s$,
 - and generates the state $(s \setminus del) \cup add$
 - (applicability of action sequences is defined as usual)



A POCL planning problem is a tuple $\mathcal{P} = \langle \mathcal{D}, P_{init} \rangle$ with:

- \mathcal{D} is the planning domain
- P_{init} is the initial plan



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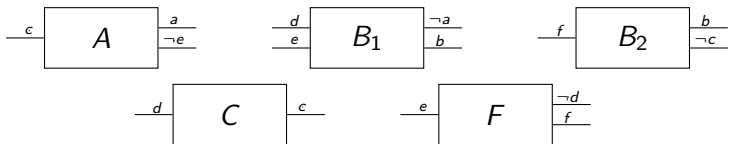
- \mathcal{D} is the planning domain
- P_{init} is the initial plan

A solution to \mathcal{P} is a plan P , s.t.:

- every precondition is supported by a causal link
- there are no causal threats



Planning domain \mathcal{D} :



Initial plan P_{init} :





flaws:

modifications:

open precondition: (a:goal)

 A

open precondition: (b:goal)

 B_1, B_2 



flaws:

modifications:

open precondition: (a:goal)

A

open precondition: (b:goal)

B_1, B_2





flaws:

modifications:

open precondition: (b:goal)

 B_1, B_2

open precondition: (c:A)

 C 



flaws:

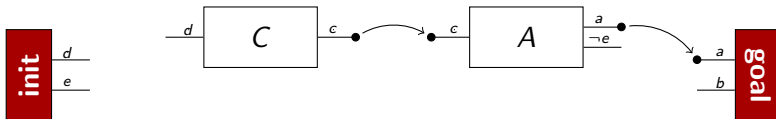
modifications:

open precondition: (b:goal)

 B_1, B_2

open precondition: (c:A)

 C 



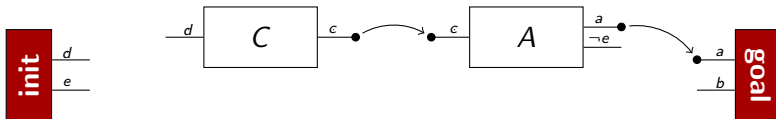
flaws:

modifications:

open precondition: (b :goal) B_1, B_2 open precondition: (d : C)

init





flaws:

modifications:

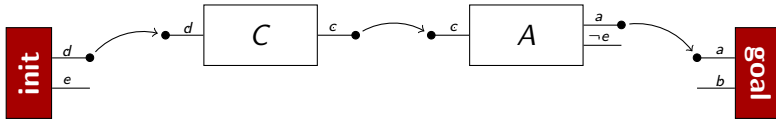
open precondition: (b:goal)

 B_1, B_2

open precondition: (d:C)

init

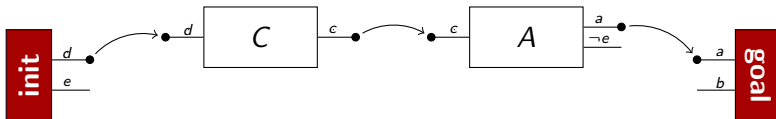




flaws:

modifications:

open precondition: $(b:\text{goal})$ B_1, B_2 

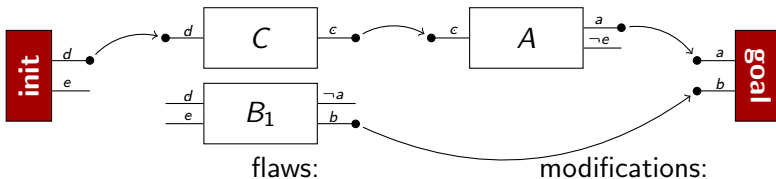


flaws:

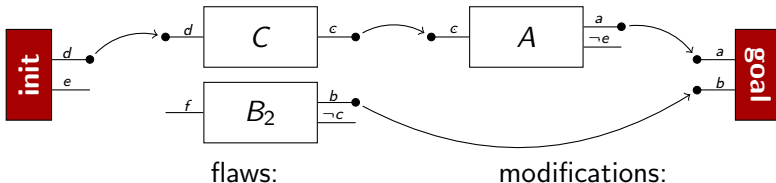
modifications:

open precondition: (b:goal)

 B_1, B_2 



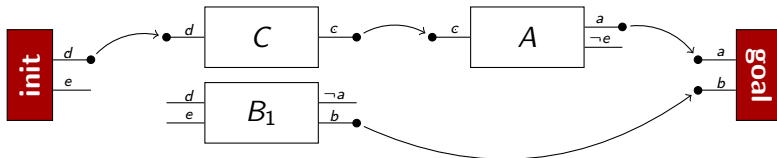
open precondition: $(d:B_1)$
open precondition: $(e:B_1)$
causal threat: $(B_1:\neg a), (A, \text{goal})$

$$\begin{array}{c} \text{init} \\ \text{init} \\ B_1 < A \end{array}$$


open precondition: $(f:B_2)$
causal threat: $(B_2:\neg c), (C,A)$

$$\overline{F}$$

$$B_2 < C, A < B_2$$

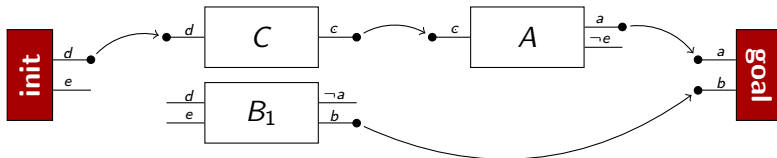
flaws:

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open precondition: $(d:B_1)$
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init
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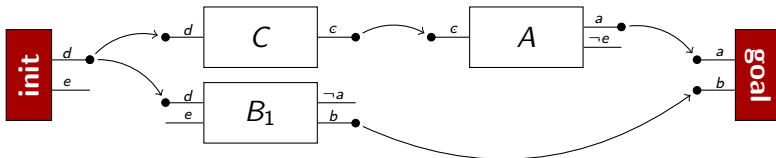
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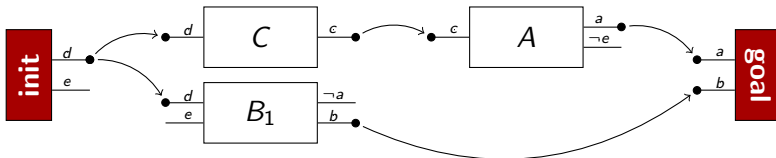


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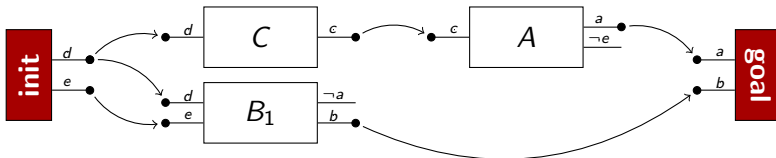


flaws:

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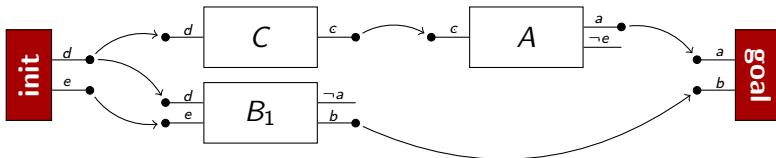
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flaws:

modifications:

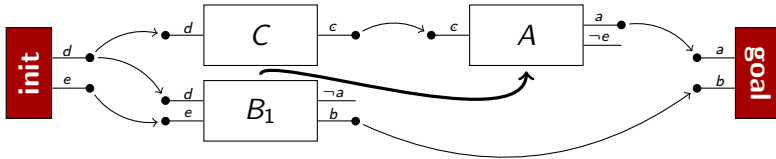
causal threat: $(B_1: \neg a), (A, \text{goal})$ $B_1 < A$ causal threat: $(A: \neg e), (\text{init}, B_1)$ $B_1 < A$ 



flaws:

modifications:

causal threat: $(B_1: \neg a), (A, \text{goal})$ $B_1 < A$ causal threat: $(A: \neg e), (\text{init}, B_1)$ $B_1 < A$ 



Choice points of search algorithm:

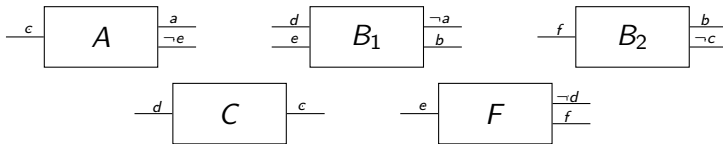
- Flaw selection like fewest modifications
- Plan selection via informed search like A^* with heuristics

In this work: Delete Relaxation Heuristic for plan selection.

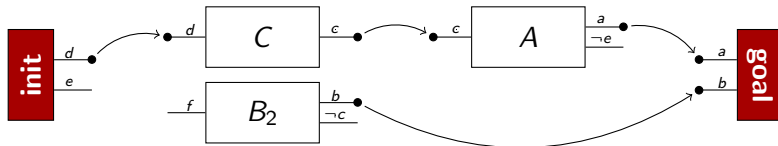
- **How hard is that estimate?** (theoretically)
- **How to estimate?** (practically)



Planning domain \mathcal{D} :



Plan to estimate:



flaws:

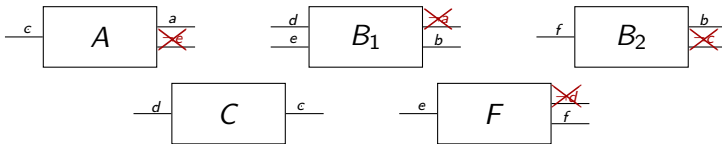
modifications:

open precondition: $(f:B_2)$
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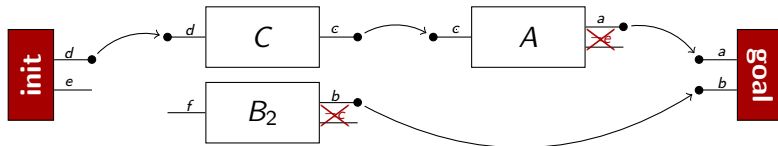
F
 $B_2 < C, A < B_2$



Planning domain \mathcal{D} :



Plan to estimate:



flaws:

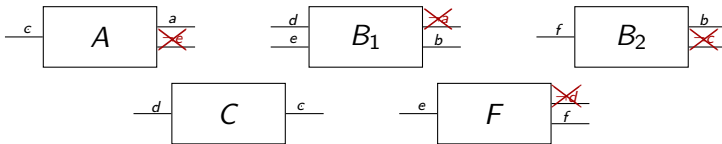
modifications:

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~~causal threat: $(B_2:\neg c), (C,A)$~~

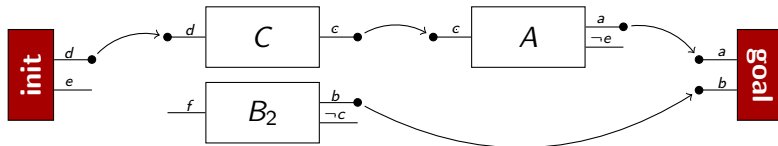
F
 ~~$B_2 \prec C, A \prec B_2$~~



Planning domain \mathcal{D} :



Plan to estimate:



flaws:

modifications:

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 $B_2 < C, A < B_2$



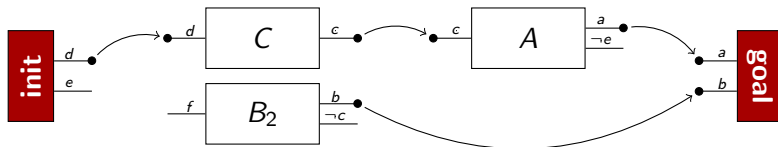
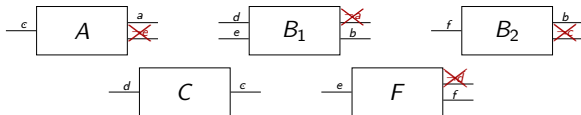
Theorem:

It is in **NP** to decide whether there is a solution for a relaxed planning domain and a non-relaxed plan.



Proof:

- Guess a non-conflicting linearization
- Apply all relaxed actions in initial state
- Apply the first plan step in the resulting state
- Continue until the last plan step (including goal) is applied



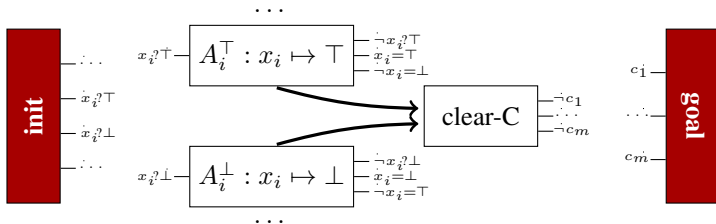
Theorem:

It is **NP-hard** to decide whether there is a solution for a relaxed planning domain and a non-relaxed plan.



Proof:

- Reduction from CNF-SAT. $X = \{x^1, \dots, x^n\}$ is a set of boolean variables. $C = \{c^1, \dots, c^m\}$ is a set of clauses; each c^j represents a disjunction of variables of X .
- Construct a delete relaxed domain and a non-relaxed plan, s.t. its solutions are isomorphic to the CNF-SAT's solutions.
- The order of A_i^\top and A_i^\perp encodes the truth assignment of $x^i \in X$. For $x^i \in c^j$, an action $C_{ij} = \langle \{x_i = \top\}, \{c_j\}, \emptyset \rangle$ encodes truth assignment of clause $c^j \in C$.



Idea:

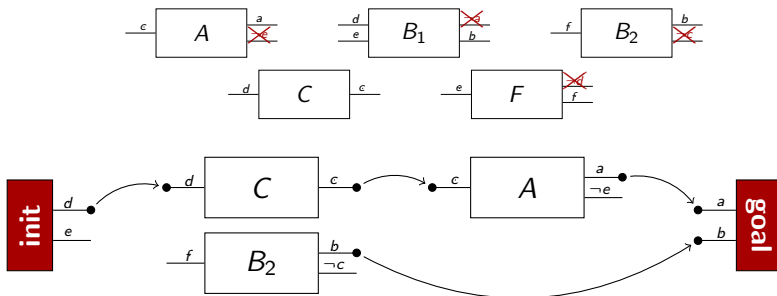
- Sample a fixed number of linearizations (this approximates the “guessing” part “)
- For each linearization: build relaxed solution as in the membership proof
- Heuristic estimate is cheapest solution



Properties:

- Does not ignore present causal links:
 - Only non-conflicting relaxed actions are used
 - Only non-conflicting linearizations are used
- Does not ignore negative effects of present plan steps

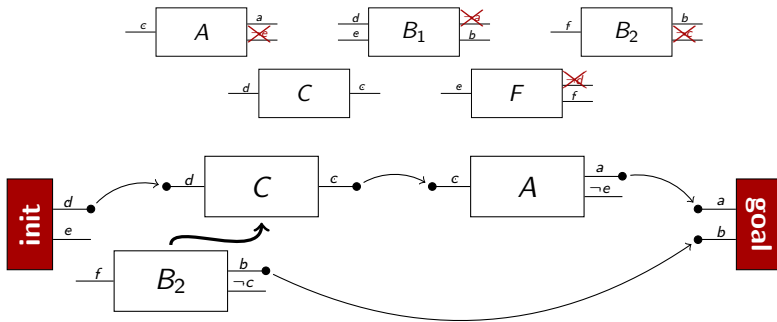
Example:



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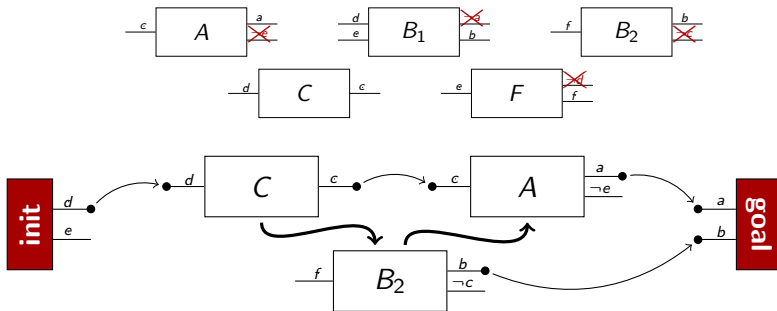
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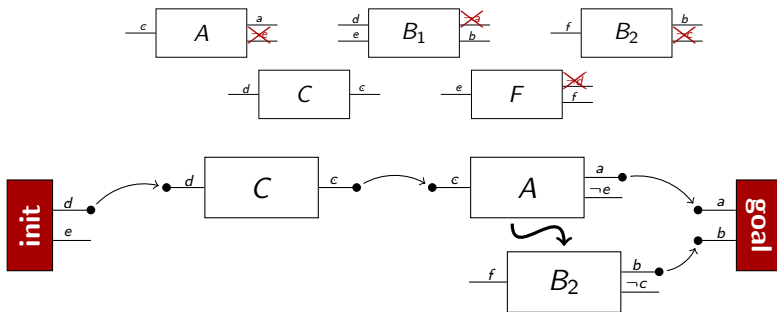
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Example:



Setting:

- Evaluation was done on 20 domains taken from the first 5 International Planning Competitions (IPCs).
- Comparison was done with:
 - Add and Relax Heuristic (for POCL planning)
 - 12 variants of SampleFF
(1, 3, 10, 30 samples, action filter feature on/off)



Results:

- number of solved problem instances overall (of 446)
Add: 292, Relax: 227/194, Sample-FF: 127 to 187
- For one problem, only Sample-FF disproved it
- Depending on the Sample-FF configuration, 1% to 46% of all plans had only invalid linearizations



Summary:

- Delete relaxation only for the domain is **NP-complete**
- We developed a new heuristic for POCL planning (Sample-FF)
 - pro: more pruning power/tighter estimates
 - con: sampled linearizations are often *all* invalid

Outlook:

- Solve the “con” problem
- More intelligent sampling. Using causal graphs?
- Lifting the heuristic



Domain	n	Add	Relax*	Relax	Sample-FF											
					front: ⊥ end: ⊥				front: ⊥ end: ⊤				front: ⊤ end: ⊤			
					1	3	10	30	1	3	10	30	1	3	10	30
grid	5	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1
gripper	20	14	20	7	1	1	1	1	1	2	1	1	2	3	3	2
logistics	20	12	8	7	8	5	6	6	6	7	6	5	0	0	1	1
movie	30	30	30	30	30	30	30	30	30	30	30	30	30	30	30	30
mystery	20	8	8	9	10	11	9	9	10	11	10	9	12	12	11	11
mystery-prime	20	3	3	3	3	4	6	5	5	4	4	4	6	6	6	6
blocks	21	4	5	7	5	5	6	6	4	3	3	3	5	3	2	0
logistics	28	28	28	27	22	23	23	24	21	19	20	21	15	13	14	15
miconic	100	100	49	39	40	40	37	35	39	41	37	32	15	16	18	20
depot	22	2	2	1	1	1	1	1	0	1	1	1	2	2	3	2
driverlog	20	7	9	7	11	9	10	9	9	10	9	8	8	7	9	7
rover	20	20	18	19	13	14	15	15	11	11	12	11	7	9	9	9
zeno-travel	10	4	5	3	3	5	4	5	4	3	4	4	1	1	1	1
airport	20	18	15	9	10	11	11	11	7	10	10	10	7	8	6	4
pipesworld-noTankage	10	8	1	2	2	3	5	3	2	1	2	1	1	4	3	5
pipesworld-Tankage	10	1	1	1	1	1	1	1	1	1	1	1	0	0	0	0
satellite	20	16	7	7	7	5	6	6	5	5	7	5	1	2	3	3
pipesworld	10	1	1	1	1	1	1	1	1	1	1	1	0	0	0	0
storage	20	7	6	4	6	7	9	8	6	7	7	6	9	9	10	10
tpp	20	19	11	11	8	7	6	7	6	6	6	6	5	6	6	6
total	446	292	227	194	182	183	187	183	168	173	171	159	127	132	136	133

Variables: $X = \{x^1, \dots, x^n\}$, Clauses: $C = \{c^1, \dots, c^m\}$

$$\mathcal{V} = \{x_i? \top, x_i? \perp, x_i = \top, x_i = \perp \mid i \in \{1, \dots, n\}\} \cup \{c_i \mid i \in \{1, \dots, m\}\}$$

$$\mathcal{A} = \{dr-A_i^\top, dr-A_i^\perp \mid i \in \{1, \dots, n\}\} \cup \{dr-clear-C\} \cup \\ \{C_{ij}^\top \mid c^j \in C, x^i \in c^j\} \cup \{C_{ij}^\perp \mid c^j \in C, \neg x^i \in c^j\}$$

$P = (PS, \prec, CL)$ and

$$PS = \{a_0, a_\infty, clear-C\} \cup \{A_i^\top, A_i^\perp \mid i \in \{1, \dots, n\}\}$$

$$\prec = \{(l_i^\top, l_c), (l_i^\perp, l_c) \mid i \in \{1, \dots, n\}\}$$

$$CL = \emptyset$$

$$a_0 = (\emptyset, \{x_i? \top, x_i? \perp \mid i \in \{1, \dots, n\}\}, \emptyset)$$

$$a_\infty = (\{c_i \mid i \in \{1, \dots, m\}\}, \emptyset, \emptyset)$$

$$A_i^\top = (\{x_i? \top\}, \{x_i = \top\}, \{x_i? \top, x_i = \perp\}), dr-A_i^\top = (\{x_i? \top\}, \{x_i = \top\}, \emptyset)$$

$$A_i^\perp = (\{x_i? \perp\}, \{x_i = \perp\}, \{x_i? \perp, x_i = \top\}), dr-A_i^\perp = (\{x_i? \perp\}, \{x_i = \perp\}, \emptyset)$$

$$clear-C = (\emptyset, \emptyset, \{c_1, \dots, c_m\}), dr-clear-C = (\emptyset, \emptyset, \emptyset)$$

$$C_{ij}^\top = (\{x_i = \top\}, \{c_j\}, \emptyset), C_{ij}^\perp = (\{x_i = \perp\}, \{c_j\}, \emptyset)$$